

## The Propositional Content of Predicate-Calculus wfs; Predicate-Calculus Tautologies

In this handout, I will construct a bijection  $\mathcal{T}$  from the set of  $\mathcal{L}$ -wfs onto the set of  $L$ -wfs, the purpose of which is to isolate the  $L$ -content of each  $\mathcal{L}$ -wf. I will then use  $\mathcal{T}$  to identify the  $\mathcal{L}$ -tautologies. My construction seems to me to be of exactly the same flavor as are others one encounters in this subject: completely precise but perhaps not completely transparent.

**Step 1.  $L$ -complexity.** Before defining  $\mathcal{T}$ , I need to construct an integer-valued function on the  $\mathcal{L}$ -wfs that will measure their propositional complexity, by counting up the number of “ $L$ -visible” occurrences of  $\sim$  and  $\rightarrow$ . I will first isolate the  $\mathcal{L}$ -wfs of propositional complexity zero; these are the  $\mathcal{L}$ -wfs in which the Propositional Calculus can discern no logical structure whatsoever.

**Definition 1.** Say that  $\mathcal{L}$ -wf  $\mathcal{A}$  is  *$L$ -irreducible* iff it is either atomic or of the form  $(\forall x_i)\mathcal{B}$  for some  $\mathcal{L}$ -wf  $\mathcal{B}$ .

I can now construct the  $L$ -complexity function  $d$  that I will need.

**Definition 2.** Define the  *$L$ -complexity*  $d$  of  $\mathcal{L}$ -wf  $\mathcal{A}$  as follows:

- $\rightarrow$ : If  $\mathcal{A}$  is  $L$ -irreducible, then  $d(\mathcal{A}) := 0$ ;
- $\rightarrow$ : If  $\mathcal{A}$  is of the form  $(\sim\mathcal{B})$ , then  $d(\mathcal{A}) := d(\mathcal{B}) + 1$ ;
- $\rightarrow$ : If  $\mathcal{A}$  is of the form  $(\mathcal{B} \rightarrow \mathcal{C})$ , then  $d(\mathcal{A}) := d(\mathcal{B}) + d(\mathcal{C}) + 1$ .

**Step 2. The bijection  $\mathcal{T}$ .** I will define a function  $\mathcal{T}(\mathcal{A}) = \hat{\mathcal{A}}$  (from  $\mathcal{L}$ -wfs to  $L$ -wfs) which will be designed to isolate the  $L$ -content of each  $\mathcal{L}$ -wf  $\mathcal{A}$ .  $\mathcal{T}$  will be defined by recursion over  $d(\mathcal{A})$ .

$\rightarrow d = 0$ :

Choose any enumeration  $(\mathcal{A}_1, \mathcal{A}_2, \mathcal{A}_3, \dots)$  of the the  $L$ -irreducible wfs and put

$$\mathcal{T}(\mathcal{A}_i) := p_i, \quad i = 1, 2, 3, \dots$$

$\rightarrow d > 0$ :

- $(\alpha)$ , if  $\mathcal{A}$  is of the form  $(\sim\mathcal{B})$ , then put  $\mathcal{T}(\mathcal{A}) := (\sim\mathcal{T}(\mathcal{B}))$ ;
- $(\beta)$ , if  $\mathcal{A}$  is of the form  $(\mathcal{B} \rightarrow \mathcal{C})$ , then put  $\mathcal{T}(\mathcal{A}) := (\mathcal{T}(\mathcal{B}) \rightarrow \mathcal{T}(\mathcal{C}))$ .

**Exercise 1.** Show that  $\mathcal{T}$  is indeed a bijection from the set of  $\mathcal{L}$ -wfs onto the set of  $L$ -wfs.

**Step 3.  $L$ -valuations from  $\mathcal{L}$ -valuations.** Let  $v$  be any  $\mathcal{L}$ -valuation in any interpretation of  $\mathcal{L}$ . Associate to  $v$  the  $L$ -valuation  $\tilde{v}$  obtained by putting, for each  $L$ -wf  $\hat{\mathcal{A}}$ ,

$$\tilde{v}(\hat{\mathcal{A}}) := f_v\left(\mathcal{T}^{-1}(\hat{\mathcal{A}})\right).$$

**Exercise 2.** Show that  $\tilde{v}$  is a propositional-calculus valuation. (That is, show that it satisfies the definition.)

**Definition 3.** Say that  $\mathcal{L}$ -wf  $\mathcal{A}$  is a *tautology in  $\mathcal{L}$*  iff  $\mathcal{T}(\mathcal{A})$  is a tautology in  $L$ .

**Exercise 3.** Show that every tautology in  $\mathcal{L}$  is logically valid.